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## FULL EDGE-FRIENDLY INDEX SETS OF COMPLETE BIPARTITE GRAPHS

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ABSTRACT. Let  $G = (V, E)$  be a simple graph. An edge labeling  $f : E \rightarrow \{0, 1\}$  induces a vertex labeling  $f^+ : V \rightarrow \mathbb{Z}_2$  defined by  $f^+(v) \equiv \sum_{uv \in E} f(uv) \pmod{2}$  for each  $v \in V$ , where  $\mathbb{Z}_2 = \{0, 1\}$  is the additive group of order 2. For  $i \in \{0, 1\}$ , let  $e_f(i) = |f^{-1}(i)|$  and  $v_f(i) = |(f^+)^{-1}(i)|$ . A labeling  $f$  is called edge-friendly if  $|e_f(1) - e_f(0)| \leq 1$ .  $I_f(G) = v_f(1) - v_f(0)$  is called the edge-friendly index of  $G$  under an edge-friendly labeling  $f$ . The full edge-friendly index set of a graph  $G$  is the set of all possible edge-friendly indices of  $G$ . Full edge-friendly index sets of complete bipartite graphs will be determined.

### 1. Introduction

Let  $G = (V, E)$  be a simple graph. An edge labeling  $f : E \rightarrow \{0, 1\} \subset \mathbb{N}$  induces a vertex labeling  $f^+ : V \rightarrow \mathbb{Z}_2$  defined by  $f^+(v) \equiv \sum_{uv \in E} f(uv) \pmod{2}$  for each  $v \in V$ , where  $\mathbb{Z}_2 = \{0, 1\}$  is the additive group of order 2. We sometimes view the value of  $f^+(v)$  as an integer. For  $i \in \{0, 1\}$ , let  $e_f(i) = |f^{-1}(i)|$  and  $v_f(i) = |(f^+)^{-1}(i)|$ . Let  $I_f(G) = v_f(1) - v_f(0)$ . An edge labeling  $f$  is *edge-friendly* if  $|e_f(1) - e_f(0)| \leq 1$ . The concept of edge-friendly index maybe first introduced by Lee and Ng [4] on considering edge cordial labeling. Unfortunately, we cannot find this paper even through we have asked the authors with response that they also do not have a reprint. Readers are referred to [1] for detail about edge cordiality.

The number  $I_f(G)$  is called the *edge-friendly index* of  $G$  under  $f$  if  $f$  is an edge-friendly labeling of  $G$ . The set  $\text{FEFI}(G) = \{I_f(G) \mid f \text{ is edge-friendly}\}$  is called the *full edge-friendly index set* of  $G$ . This is a dual concept of full friendly index set which was first introduced by the author and H. Kwong [10]. Readers who are interested on friendly index or friendly index set may refer to [2, 3, 5, 6, 8–16].

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In [7], the author proposed a conjecture that

**Conjecture 1.1.**

$$\text{FEFI}(K_{m,n}) = \begin{cases} \{4j - (m+n) \mid 1 \leq j \leq \lfloor (m+n)/2 \rfloor\}, & \text{if } n \equiv 2 \pmod{4} \text{ and } m = 2 \text{ or } m \text{ is odd;} \\ \{4j - (m+n) \mid 1 \leq j \leq \lfloor (m+n)/2 \rfloor\}, & \text{if } m \equiv 2 \pmod{4} \text{ and } n = 2 \text{ or } n \text{ is odd;} \\ \{4j - (m+n) \mid 0 \leq j \leq \lfloor (m+n)/2 \rfloor\}, & \text{otherwise.} \end{cases}$$

This paper is a continuation of [7]. We shall determine full edge-friendly index sets of complete bipartite graphs  $K_{m,n}$  and settle the above conjecture.

## 2. Some Basic Properties

In the following, all graphs are simple and connected. The codomain of any edge labeling is  $\mathbb{Z}_2$ . Suppose  $f$  is an edge labeling. A vertex (resp. an edge) is called an  $i$ -vertex (resp.  $i$ -edge) under  $f$  if it is labeled by  $i \in \{0, 1\}$ . Notation and concepts not defined here are referred to [17].

Suppose  $G$  is a graph of order  $p$ . Since  $v_f(1) + v_f(0) = p$  for any edge labeling  $f$  of  $G$ ,  $I_f(G) = 2v_f(1) - p$ . Thus, it suffices to study the number of 1-vertices instead of studying the edge-friendly index of  $G$  under  $f$ .

**Lemma 2.1** ([4, 7]). *Let  $f$  be any edge labeling of a graph  $G = (V, E)$ . Then  $v_f(1)$  must be even.*

By means of the above lemma, we may write  $v_f(1) = 2j$  for some  $j$  with  $0 \leq j \leq \lfloor p/2 \rfloor$ , where  $f$  is an edge labeling of a graph  $G$  of order  $p$ . So  $I_f(G) = 4j - p$  for some  $j$ ,  $0 \leq j \leq \lfloor p/2 \rfloor$ . It implies that

$$\text{FEFI}(G) \subseteq \{4j - p \mid 0 \leq j \leq \lfloor p/2 \rfloor\}.$$

A labeling matrix  $L_f(G)$  for an edge labeling  $f$  of a graph  $G$  is a matrix whose rows and columns are indexed by the vertices of  $G$  and the  $(u, v)$ -entry is  $f(uv)$  if  $uv \in E$ , and is  $*$  otherwise.

Suppose  $L_f(G)$  is a labeling matrix for the edge labeling  $f$  of  $G$ . If we view the entries of  $L_f(G)$  as elements in  $\mathbb{Z}_2$ , then  $f^+(v)$  is the  $v$ -row sum (as well as  $v$ -column sum), where entries with  $*$  will be treated as 0.

Let  $X = \{x_1, \dots, x_m\}$  and  $Y = \{y_1, \dots, y_n\}$  be the bipartition of the complete bipartite graph  $K_{m,n}$ . Under this indexing of vertices, a labeling matrix for any edge labeling  $f$  is of the form

$$\begin{pmatrix} \star_m & A \\ A^T & \star_n \end{pmatrix},$$

where  $\star_r$  is a square matrix of order  $r$  with all entries being  $*$  and  $A$  is an  $m \times n$  matrix whose entries are elements of  $\mathbb{Z}_2$ . So the multi-set of row sums and column sums of  $A$  is equal to the sequence  $\{f^+(x_1), \dots, f^+(x_m), f^+(y_1), \dots, f^+(y_n)\}$ . Thus, we shall only consider such matrix  $A$  and we shall denote it as  $A_f(G)$  when there is some ambiguity. Thus, we shall use such matrix  $A_f(G)$  (or  $A$ ) to define an edge labeling  $f$ . Let  $v_A(1)$  denote the number of 1's being row sum or column sum. Then  $v_A(1) = v_f(1)$ . Similarly, we may define  $v_A(0)$ , which will equal to  $v_f(0)$ . Also we may define  $e_A(1)$  and  $e_A(0)$  to be the number of 1 and 0 used to form the matrix  $A$ , respectively. So  $e_A(i) = e_f(i)$ ,  $i = 0, 1$ .

An  $m \times n$  matrix  $A$  satisfying the following conditions is called a *friendly matrix* of  $K_{m,n}$ :

1. Each entry of  $A$  is either 1 or 0;
2.  $|e_A(1) - e_A(0)| \leq 1$ .

Actually, in Conjecture 1.1,  $2j$  is equal to  $v_A(1)$  for some friendly matrix  $A$ . Since we only consider the value of  $v_A(1)$  later, we simple write this value as  $s(A)$  and called it the  $s$ -value of  $A$ .

It was listed in [7] that

$$\text{FEFI}(K_{1,n}) = \begin{cases} \{-2, 2\}, & n = 4k + 1; \\ \{1\}, & n = 4k + 2; \\ \{0\}, & n = 4k + 3; \\ \{-1\}, & n = 4k + 4, \end{cases}$$

where  $k \geq 0$ .

In the following sections, we want to find some friendly matrices  $A$  of  $K_{m,n}$  such that  $v_A(1)$  run through all the possible  $s$ -values, where  $m, n \geq 2$ .

### 3. Full Edge-friendly Index Sets of $K_{2,n}$

It is known from [7, Example 4.5] that Conjecture 1.1 holds for  $n \equiv 2 \pmod{4}$ . So we only need to deal with  $n = 2k + 1$  or  $n = 4k$  for  $k \geq 1$ .

For easy to describe some matrices, let  $J_{m,n}$  be the  $m \times n$  matrix whose entries are 1 and  $O_{m,n}$  be the  $m \times n$  zero matrix.

We first consider  $n = 2k + 1$ , for some  $k \geq 1$ . We want to show that

$$(3.1) \quad \text{FEFI}(K_{2,2k+1}) = \{4j - 2k - 3 \mid 1 \leq j \leq k + 1\}$$

Let the block matrix  $A_1 = \begin{pmatrix} J_{2,k} & O_{2,k} & 1 \\ & & 0 \end{pmatrix}$  which is a friendly matrix of  $K_{2,2k+1}$ . Clearly  $s(A_1) = 2$ .

For  $1 \leq i \leq k$ , let  $A_{i+1}$  be the matrix obtained from  $A_i$  by swapping  $(A_i)_{1,i}$  (the  $(1, i)$ -entry of  $A_i$ ) with  $(A_i)_{1,k+i}$ . Then  $s(A_{i+1}) = s(A_i) + 2 = 2i + 2$ . Hence we obtain each even number between 2 and  $2(k + 1)$  as a value of  $s(A)$  for some friendly matrix  $A$ . So we get (3.1).

Next, we consider  $n = 4k$ , for some  $k \geq 1$ . Let the block matrix  $B_0 = \begin{pmatrix} J_{2,2k} & O_{2,2k} \end{pmatrix}$  which is a friendly matrix of  $K_{2,4k}$ . Clearly  $s(B_0) = 0$ . By a similar procedure as above, we will get

$$\{4j - 4k - 2 \mid 0 \leq j \leq 2k\} \subseteq \text{FEFI}(K_{2,4k})$$

Following lemma was proved at [7, Lemma 4.2]:

**Lemma 3.1.** *Suppose  $m$  and  $n$  are even. There is a friendly matrix  $M$  of  $K_{m,n}$  such that  $v_M(1) = m+n$ .*

Combining Lemma 3.1 and the above discussion, we have

$$\text{FEFI}(K_{2,4k}) = \{4j - 4k - 2 \mid 0 \leq j \leq 2k + 1\}$$

So we have

**Theorem 3.2.** *Conjecture 1.1 holds when  $m = 2$ .*

For now on, we assume  $m, n \geq 3$ .

4. Full Edge-friendly Index Sets of  $K_{m,n}$  with even  $m$

We list some useful matrices which were defined in [7].

$$A_{2s,4} = \begin{pmatrix} J_{2s,2} & O_{2s,2} \end{pmatrix} \text{ for } s \geq 1, \quad A_{3,4} = \begin{pmatrix} 1 & 1 & 0 & 0 \\ 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 0 \end{pmatrix},$$

$$D_s = \begin{pmatrix} J_{s,6} \\ O_{s,6} \end{pmatrix} \text{ for } s \geq 1, \quad A_{6,6} = \begin{pmatrix} 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 & 0 & 0 \\ 1 & 1 & 1 & 0 & 1 & 0 \\ 1 & 1 & 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \end{pmatrix}.$$

(4.1)  $A_{2s,4k} = J_{1,k} \otimes A_{2s,4}$ , the Kronecker product of  $J_{1,k}$  and  $A_{2s,4}$ ,

(4.2)  $A_{2s+1,4k} = J_{1,k} \otimes \begin{pmatrix} A_{2s-2,4} \\ A_{3,4} \end{pmatrix}$ ,

(4.3)  $A_{4h+2,4k+2} = \left( \begin{array}{c|c} J_{1,k-1} \otimes A_{4h-4,4} & D_{2h-2} \\ \hline J_{2,4k-4} \otimes A_{3,4} & A_{6,6} \end{array} \right)$

Before finding the required friendly matrices, we define some procedures:

**Procedure R:** Let  $R_0$  be a given  $m \times 2t$  friendly matrix. For  $1 \leq i \leq t$ , let  $R_i$  be the matrix obtained from  $R_{i-1}$  by swapping  $(R_{i-1})_{1,i}$  with  $(R_{i-1})_{1,t+i}$ .

**Procedure C:** Let  $C_0$  be a given  $2s \times n$  friendly matrix. For  $1 \leq i \leq s$ , let  $C_i$  be the matrix obtained from  $C_{i-1}$  by swapping  $(C_{i-1})_{i,1}$  with  $(C_{i-1})_{s+i,1}$ .

We first consider  $m = 4h + 2$  with  $h \geq 1$ .

**Case 1.1:** Suppose  $n = 4k, k \geq 1$ . In this case, we want to find a friendly matrix  $A$  such that  $s(A) = 2j$  for each  $j$ , where  $0 \leq j \leq 2h + 2k + 1$ .

Let  $B_0 = \begin{pmatrix} J_{4h+2,2k} & O_{4h+2,2k} \end{pmatrix}$ . Then  $s(B_0) = 0$ . Applying Procedure R to  $B_0$ , we get  $B_i$ , for  $1 \leq i \leq 2k$ . It is easy to see that  $s(B_i) = 2i$ .

Let  $C_0 = \begin{pmatrix} J_{2h+1,4k} \\ O_{2h+1,4k} \end{pmatrix}$ . Then  $s(C_0) = 4k$ . Applying Procedure C to  $C_0$ , we get  $C_i$  for  $1 \leq i \leq 2h + 1$ . Clearly  $s(C_i) = 4k + 2i$ .

Hence we get the result.

**Example 4.1.** Consider the graph  $K_{6,8}$ .

Let  $B_0 = \begin{pmatrix} 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 \end{pmatrix}$  and  $C_0 = \begin{pmatrix} 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{pmatrix}$ .



Step 1: We have

$$B_0 \rightarrow B_1 = \left( \begin{array}{cccc|cccc} 1 & 1 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 0 & 1 & 1 & 1 & 1 & 0 & 0 \\ 1 & 0 & 1 & 0 & 1 & 1 & 1 & 0 & 1 & 0 \\ 0 & 1 & 1 & 0 & 1 & 1 & 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{array} \right) \rightarrow B_2 = \left( \begin{array}{cccc|cccc} 1 & 1 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 0 & 1 & 1 & 1 & 1 & 0 & 0 \\ 1 & 0 & 1 & 0 & 1 & 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{array} \right)$$

Step 2: We have

$$B_2 \rightarrow \left( \begin{array}{cccc|cccc} 1 & 1 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 0 & 1 & 1 & 1 & 1 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{array} \right) \rightarrow \left( \begin{array}{cccc|cccc} 1 & 1 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 0 & 1 & 1 & 1 & 1 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \end{array} \right)$$

Then the corresponding  $s$ -values of these matrices are 0, 2, 4, 6, 8. After applying Procedure C to  $C_0$ , we obtain the  $s$ -values being 10, 12, 14, 16. The last matrix of this step is

$$C_3 = \begin{pmatrix} 0 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{pmatrix}.$$

**Case 1.3:** Suppose  $n = 2t + 1 \geq 4h + 2$ . In this case, we want to find a friendly matrix  $A$  such that  $s(A) = 2j$  for each  $j$ , where  $1 \leq j \leq 2h + t + 1$ .

To make the presentation easier to follow, we consider the graph  $K_{2t+1,4h+2}$ , which is isomorphic to  $K_{4h+2,2t+1}$ .

Let  $Z_{2t+1,2} = \begin{pmatrix} J_{t,2} \\ 1 \ 0 \\ O_{t,2} \end{pmatrix}$  and  $B_1 = A_{2t+1,4h+2} = (A_{2t+1,4h} \ Z_{2t+1,2})$ , where  $A_{2t+1,4h}$  is defined in (4.2). It is known that  $s(B_1) = 2$  (c.f. [7]).

Do the same procedure as Step 1 of Case 1.2, we get  $2h$  matrices whose  $s$ -values run through the even numbers between 4 and  $4h + 2$ . After performing this step, let the last matrix be  $B$ . Note that the submatrix consisting of the last two columns of  $B$  is still the matrix  $Z_{2t+1,2}$ . For  $1 \leq i \leq t$ , swap  $(Z_{2t+1,2})_{i,1}$  with  $(Z_{2t+1,2})_{t+1+i,1}$  in the matrix  $B$ . Then we obtain  $t$  matrices whose  $s$ -values run through the even numbers between  $4h + 4$  and  $4h + 2 + 2t$ .

Hence we get the result.

**Example 4.3.** Consider the graph  $K_{6,7}$ . From the above discussion we consider the graph  $K_{7,6}$  instead.

$$\text{Let } B_1 = \left( \begin{array}{cccc|cc} 1 & 1 & 0 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 & 1 & 0 \\ 1 & 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 0 \end{array} \right).$$

Applying the same procedure as Step 1 of Case 1.2, we have

$$B_1 \rightarrow B_2 = \left( \begin{array}{cccc|cc} 1 & 1 & 0 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 & 1 & 1 \\ \hline 1 & 1 & 0 & 0 & 1 & 0 \\ 0 & 1 & 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 0 \end{array} \right) \rightarrow B_3 = \left( \begin{array}{cccc|cc} 1 & 1 & 0 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 & 1 & 1 \\ \hline 1 & 1 & 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 & 0 & 0 \\ 0 & 1 & 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 0 \end{array} \right)$$

Then the corresponding  $s$ -values of these matrices are 2, 4, 6. Swapping entries of the submatrix  $Z_{7,2}$ , we have

$$B_3 \rightarrow \left( \begin{array}{cccc|cc} 1 & 1 & 0 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 & 1 & 1 \\ \hline 1 & 1 & 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 & 1 & 0 \\ 0 & 1 & 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 0 \end{array} \right) \rightarrow \left( \begin{array}{cccc|cc} 1 & 1 & 0 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 & 1 & 1 \\ \hline 1 & 1 & 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 & 1 & 0 \\ 0 & 1 & 1 & 0 & 1 & 0 \\ 1 & 0 & 1 & 0 & 0 & 0 \end{array} \right) \rightarrow \left( \begin{array}{cccc|cc} 1 & 1 & 0 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 & 0 & 1 \\ \hline 1 & 1 & 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 & 1 & 0 \\ 0 & 1 & 1 & 0 & 1 & 0 \\ 1 & 0 & 1 & 0 & 1 & 0 \end{array} \right)$$

Then the corresponding  $s$ -values of these matrices are 6, 8, 10, 12.

Next, we consider  $m = 4h$  with  $h \geq 1$ . For easy to present, we consider  $K_{n,4h}$  instead of  $K_{4h,n}$ . If  $n = 4k + 2$ , then we can refer to Case 1.1. So we only consider  $n = 4k$  and  $n = 2t + 1$ .

**Case 2.1:** Suppose  $n = 4k$ ,  $k \geq 1$ . In this case, we want to find a friendly matrix  $A$  such that  $s(A) = 2j$  for each  $j$ , where  $0 \leq j \leq 2h + 2k$ .

Let  $B_0 = \begin{pmatrix} J_{4k,2h} & O_{4k,2h} \end{pmatrix}$ . Similar to Case 1.1 we obtain matrix  $B_i$  such that  $s(B_i) = 2i$  for  $0 \leq i \leq 2h$ .

Let  $C_0 = \begin{pmatrix} J_{2k+1,2h} & O_{2k+1,2h} \\ O_{2k-1,2h} & J_{2k-1,2h} \end{pmatrix}$ . Clearly  $s(C_0) = 4h$ . Applying Procedure C to  $C_0$  (the first step is redundant), we obtain  $2k$  matrices whose  $s$ -values run through the even numbers between  $4h$  and  $4h + 4k - 2$ . Combining with the maximum value obtained in [7, Lemma 4.2], we have the result.

**Case 2.2:** Suppose  $n = 2t + 1$ ,  $t \geq 1$ . In this case, we want to find a friendly matrix  $A$  such that  $s(A) = 2j$  for each  $j$ , where  $0 \leq j \leq 2h + t$ .

Let  $B_0 = A_{2t+1,4h}$ . It is known [7] that  $s(B_0) = 0$ . Apply the procedure similar to Step 1 of Case 1.2 we obtain  $2h$  matrices whose  $s$ -values run through the even numbers between 2 and  $4h$ . The last matrix

$$B_{2h} \text{ is } J_{1,h} \otimes \begin{pmatrix} A_{2t-2,4} \\ B_{3,4} \end{pmatrix}, \text{ where } B_{3,4} = \begin{pmatrix} 0 & 0 & 1 & 1 \\ 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 0 \end{pmatrix}.$$

Before we continue the construction, we define two more procedures.

**Procedure S1:** Consider the matrix  $A_{4,4}$ . We perform the following two steps:

$$A_{4,4} = \begin{pmatrix} 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \end{pmatrix} \rightarrow A_{4,4}^{(1)} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 \\ 1 & 1 & 1 & 0 \\ 1 & 1 & 0 & 0 \end{pmatrix} \rightarrow A_{4,4}^{(2)} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 1 & 1 & 1 & 0 \\ 1 & 1 & 1 & 0 \end{pmatrix}.$$

Clearly,  $s(A_{4,4}^{(1)}) = 2$  and  $s(A_{4,4}^{(2)}) = 4$ .

**Procedure S2:** Consider the matrix  $S = \begin{pmatrix} 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 \\ 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 0 \end{pmatrix}$ . We perform the following two steps:

$$S \rightarrow S_1 = \begin{pmatrix} 0 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 0 & 1 & 1 \\ 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 0 \end{pmatrix} \rightarrow S_2 = \begin{pmatrix} 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 1 & 0 \\ 1 & 0 & 1 & 0 \end{pmatrix}$$

Clearly,  $s(S) = 4$ ,  $s(S_1) = 6$  and  $s(S_2) = 8$ .

Now we return to consider Case 2.2.

Suppose  $t = 2k + 1$ . Then the first 4 columns of  $B_{2h}$  is  $\begin{pmatrix} J_{k,1} \otimes A_{4,4} \\ B_{3,4} \end{pmatrix}$ . Applying Procedure S1 to  $A_{4,4}$  of the first 4 columns of  $B_{2h}$  one by one, we obtain  $2k$  matrices whose  $s$ -values run through the even numbers between  $4h + 2$  and  $4h + 4k$ . Combining with the maximum value obtained in [7, Lemma 4.2], we have the result.

Suppose  $t = 2k$ . Then the first 4 columns of  $B_{2h}$  is  $\begin{pmatrix} J_{k-1,1} \otimes A_{4,4} \\ S \end{pmatrix}$ . Applying Procedure S1 to  $A_{4,4}$  of the first 4 columns of  $B_{2h}$  one by one, we obtain  $2k - 2$  matrices whose  $s$ -values run through the even numbers between  $4h + 2$  and  $4h + 4k - 4$ . After that, applying Procedure S2 to  $S$  of the first 4 columns of  $B_{2h}$  we obtain two matrices whose  $s$ -values are  $4h + 4k - 2$  and  $4h + 4k$ . So we have the result.

**Example 4.4.** Consider the graph  $K_{9,4}$ . Applying a similar procedure as Step 1 of Case 1.2, Procedure S1 and then Procedure S2, we have

$$B_0 = \begin{pmatrix} 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 0 \end{pmatrix} \rightarrow B_1 = \begin{pmatrix} 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 1 & 0 \\ 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 0 \end{pmatrix} \rightarrow B_2 = \begin{pmatrix} 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 \\ 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 0 \end{pmatrix} \rightarrow \begin{pmatrix} 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 \\ 1 & 1 & 1 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 \\ 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 0 \end{pmatrix}$$

$$\rightarrow \begin{pmatrix} 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 \\ 1 & 1 & 1 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 \\ 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 0 \end{pmatrix} \rightarrow \begin{pmatrix} 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 1 & 1 & 1 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 0 \end{pmatrix} \rightarrow \begin{pmatrix} 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 1 & 1 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 1 & 1 & 1 & 0 \\ 1 & 0 & 1 & 0 \end{pmatrix}$$

Hence the corresponding  $s$ -values of these matrices are 0, 2, 4, 6, 8, 10, 12.

Combining the discussions above, we have

**Theorem 4.1.** Conjecture 1.1 holds when  $m$  is even.



5. Full Edge-friendly Index Sets of  $K_{m,n}$  with odd  $m$  and  $n$

Now, by symmetry we have to deal with three cases: (a)  $m = 4h + 3$  and  $n = 4k + 3$ ; (b)  $m = 4h + 1$  and  $n = 4k + 3$ ; (c)  $m = 4h + 1$  and  $n = 4k + 1$ .

$$\text{Let } A_{3,3} = \begin{pmatrix} 1 & 1 & 0 \\ 1 & 1 & 0 \\ 0 & 0 & 0 \end{pmatrix}, A_{4,3} = \begin{pmatrix} 1 & 0 & 1 \\ 1 & 1 & 0 \\ 0 & 1 & 1 \\ 0 & 0 & 0 \end{pmatrix}, A_{5,3} = \begin{pmatrix} 1 & 1 & 0 \\ 1 & 1 & 0 \\ 1 & 1 & 0 \\ 1 & 1 & 0 \\ 0 & 0 & 0 \end{pmatrix}, A_{4,5} = \begin{pmatrix} 1 & 1 & 1 & 1 & 0 \\ 1 & 1 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 \end{pmatrix}, \text{ and}$$

$$A_{5,5} = \begin{pmatrix} 1 & 1 & 1 & 1 & 0 \\ 1 & 1 & 1 & 1 & 0 \\ 1 & 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{pmatrix}. \text{ Note that all } s\text{-values of these friendly matrices are 0.}$$

**Case (a):** Suppose  $m = 4h + 3$  and  $m = 4k + 3$ . We start from the friendly matrix

$$A_{4h+3,4k+3} = \begin{pmatrix} A_{4h,4k} & J_{h,1} \otimes A_{4,3} \\ J_{1,k} \otimes A_{3,4} & A_{3,3} \end{pmatrix},$$

whose  $s$ -value is 0. We apply a similar Procedure R to each submatrix  $A_{3,4}$  lying in the last row of the block matrix  $A_{4h+3,4k+3}$  one by one. Then we obtain  $2k$  matrices whose  $s$ -values run through the even numbers between 2 to  $4k$ . After that, we apply Procedure C to submatrices  $A_{4,3}$  lying in the last column of the block matrix  $A_{4h+3,4k+3}$  one by one. Then we obtain  $2h$  matrices whose  $s$ -values run through the even numbers between  $2 + 4k$  to  $4h + 4k$ .

For the  $A_{3,3}$  lying at the lower-right corner of the block matrix  $A_{4h+3,4k+3}$ , we apply the following procedure:

$$A_{3,3} \rightarrow \begin{pmatrix} 1 & 0 & 1 \\ 1 & 1 & 0 \\ 0 & 0 & 0 \end{pmatrix} \rightarrow \begin{pmatrix} 1 & 1 & 1 \\ 1 & 0 & 0 \\ 0 & 0 & 0 \end{pmatrix} \rightarrow \begin{pmatrix} 1 & 1 & 1 \\ 1 & 0 & 0 \\ 1 & 0 & 0 \end{pmatrix}.$$

Note that, the last step is to replace a 0 to 1. The resulting matrix is still friendly. So we obtain three more matrices whose  $s$ -values are  $4h + 4k + 2$ ,  $4h + 4k + 4$  and  $4h + 4k + 6$ . Hence we get the result.

**Case (b):** Suppose  $m = 4h + 1$  and  $n = 4k + 3$ . We start from the friendly matrix

$$A_{4h+1,4k+3} = \left( A_{4h+1,4k} \left| \begin{array}{c} J_{h-1,1} \otimes A_{4,3} \\ A_{5,3} \end{array} \right. \right),$$

where  $A_{4h+1,4k}$  was defined in (4.2). Similar to Case (a), we apply Procedure R and Procedure C to each submatrices  $A_{3,4}$  and  $A_{4,3}$ , respectively. Then we obtain  $2k + 2h - 2$  matrices whose  $s$ -values run through the even numbers from 2 to  $4h + 4k - 4$ . After that, replace the lower right corner  $A_{5,3}$  by the following 4 matrices we will get 4 matrices whose  $s$ -values are  $4h + 4k - 2$ ,  $4h + 4k$ ,  $4h + 4k + 2$  and

$4h + 4k + 4$ :

$$\begin{pmatrix} 1 & 0 & 1 \\ 1 & 1 & 0 \\ 1 & 1 & 0 \\ 1 & 1 & 0 \\ 0 & 0 & 0 \end{pmatrix}, \begin{pmatrix} 1 & 1 & 1 \\ 1 & 0 & 0 \\ 1 & 1 & 0 \\ 1 & 1 & 0 \\ 0 & 0 & 0 \end{pmatrix}, \begin{pmatrix} 1 & 1 & 1 \\ 1 & 0 & 0 \\ 1 & 0 & 0 \\ 1 & 1 & 0 \\ 0 & 1 & 0 \end{pmatrix}, \begin{pmatrix} 1 & 1 & 1 \\ 1 & 0 & 0 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 1 & 0 \end{pmatrix}.$$

Hence we get the result.

**Case (c):** Suppose  $m = 4h + 1$  and  $m = 4k + 1$ . We start from the friendly matrix

$$A_{4h+1,4k+1} = \left( A_{4h+1,4k-4} \mid \begin{array}{c} J_{h-1,1} \otimes A_{4,5} \\ A_{5,5} \end{array} \right).$$

Similar to Case (a), we apply Procedure R and Procedure C to each submatrices  $A_{3,4}$  and  $A_{4,5}$ , respectively. Then we obtain  $2k + 2h - 4$  matrices whose  $s$ -values run through the even numbers from 2 to  $4h + 4k - 8$ . After that, replace the lower right corner  $A_{5,5}$  by the following 5 matrices we will get 5 matrices whose  $s$ -values are  $4h + 4k - 6$ ,  $4h + 4k - 4$ ,  $4h + 4k - 2$ ,  $4h + 4k$ , and  $4h + 4k + 2$ :

$$\begin{pmatrix} 1 & 1 & 1 & 0 & 0 \\ 1 & 1 & 1 & 1 & 0 \\ 1 & 1 & 0 & 1 & 0 \\ 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{pmatrix}, \begin{pmatrix} 1 & 1 & 1 & 0 & 0 \\ 1 & 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 1 & 0 \\ 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{pmatrix}, \begin{pmatrix} 0 & 1 & 1 & 1 & 0 \\ 1 & 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 1 & 0 \\ 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{pmatrix}, \begin{pmatrix} 0 & 0 & 1 & 1 & 1 \\ 1 & 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 1 & 0 \\ 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{pmatrix}, \text{ and } \begin{pmatrix} 0 & 0 & 1 & 1 & 1 \\ 1 & 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 1 & 0 \\ 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 0 \end{pmatrix}.$$

Hence we get the result.

Combining the discussions above, we have

**Theorem 5.1.** *Conjecture 1.1 holds when both  $m$  and  $n$  are odd.*

That means Conjecture 1.1 holds for any case.

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